

Import Networks, Fuzzy Relations and Semirings

Roland Glück



Universität Augsburg

April 2008
Frauenwörth

About

- algebraic approach to networks
- using fuzzy relations
- Kleene algebra

Previous Work

- Yasuo Kawahara: On the Cardinality of Relations, 2006
- Roland Glück: Networks, Semirings and Fuzzy Relations, 2007

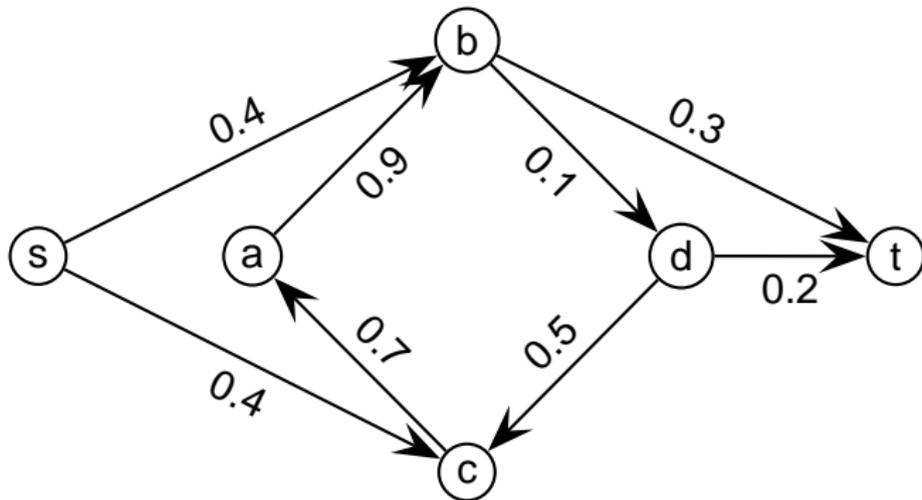
Basics

Definition

A *fuzzy relation* $\alpha : X \leftrightarrow Y$ between a set X and a set Y is a function $\alpha : X \times Y \rightarrow [0, 1]$

- natural generalisation of traditional relations
- total function!
- notation $Rel(X, Y)$ for set of all fuzzy relations between X and Y

Visualisation of Fuzzy Relations



Special Fuzzy Relations

- *empty relation* $0_{XY} : X \leftrightarrow Y$ with $0_{XY}(x, y) = 0$
- *universal relation* $\nabla_{XY} : X \leftrightarrow Y$ with $\nabla_{XY}(x, y) = 1$
- *identity relation* $id_X : X \leftrightarrow X$ with $id_X(x, y) = \delta_{xy}$

Operations on Fuzzy Relations

Let $\alpha : X \leftrightarrow Y$, $\beta : X \leftrightarrow Y$ be fuzzy relations. Then we define:

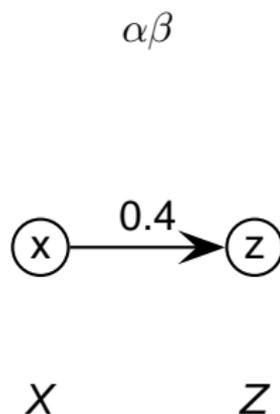
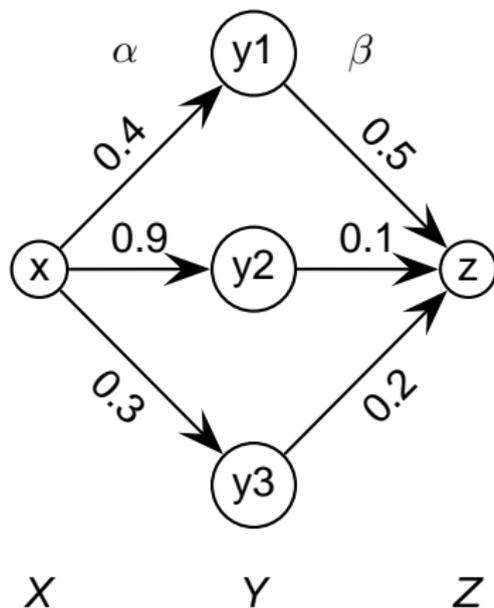
- the *join* $\alpha \sqcup \beta : X \leftrightarrow Y$ by
$$(\alpha \sqcup \beta)(x, y) := \max\{\alpha(x, y), \beta(x, y)\}$$
- the *meet* $\alpha \sqcap \beta : X \leftrightarrow Y$ by
$$(\alpha \sqcap \beta)(x, y) := \min\{\alpha(x, y), \beta(x, y)\}$$
- the *converse* $\alpha^\# : Y \leftrightarrow X$ by
$$\alpha^\#(y, x) := \alpha(x, y)$$
- the *scalar multiplication* $k \cdot \alpha$ for $k \in [0, 1]$ by
$$k \cdot \alpha(x, y) := k \cdot \alpha(x, y)$$

Composition of Fuzzy Relations

Let $\alpha : X \leftrightarrow Y$, $\beta : Y \leftrightarrow Z$ be fuzzy relations. Then we define:

- the *composition* $\alpha \circ \beta : X \leftrightarrow Z$ by
$$\alpha \circ \beta(x, z) = \max_{y \in Y} \min\{\alpha(x, y), \beta(y, z)\}$$
- abbreviation $\alpha\beta$ for $\alpha \circ \beta$
- maximally possible amount from x to z via an element of Y

Example Composition



Ordering

Lemma

The relation \sqsubseteq on $\text{Rel}(X, Y)$, defined by

$$\alpha \sqsubseteq \beta \Leftrightarrow \alpha(x, y) \leq \beta(x, y) \quad \forall (x, y) \in X \times Y$$

is a partial order on $\text{Rel}(X, Y)$.

- Note that $\alpha \sqsubseteq \beta \Leftrightarrow \alpha \sqcup \beta = \beta$

Surprise, Surprise!

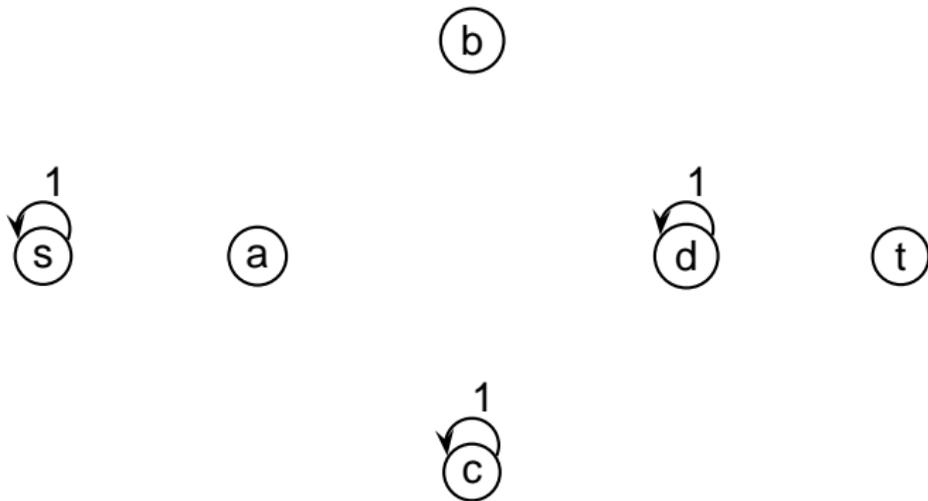
Theorem

The fuzzy endorelations over an arbitrary set X form an idempotent semiring with join as addition, composition as multiplication, 0_{XX} as zero and id_X as one. Its natural order is \sqsubseteq . and it can be extended to a Kleene algebra.

Test Relations

- A *test relation* on X is a subrelation of the identity relation id_X with a range contained in $\{0, 1\}$.
- 1:1-correspondence between test relations on X and subsets of X
- A *point relation* is a test relation corresponding to a singleton subset.
- They can be used to restrict fuzzy relations to edges beginning or ending in subsets of X .

Example Test Relation



Definition of Cardinality

Definition

The *cardinality* $|\alpha|$ of a fuzzy relation $\alpha : X \leftrightarrow Y$ is defined by

$$|\alpha| = \sum_{(x,y) \in X \times Y} \alpha(x, y)$$

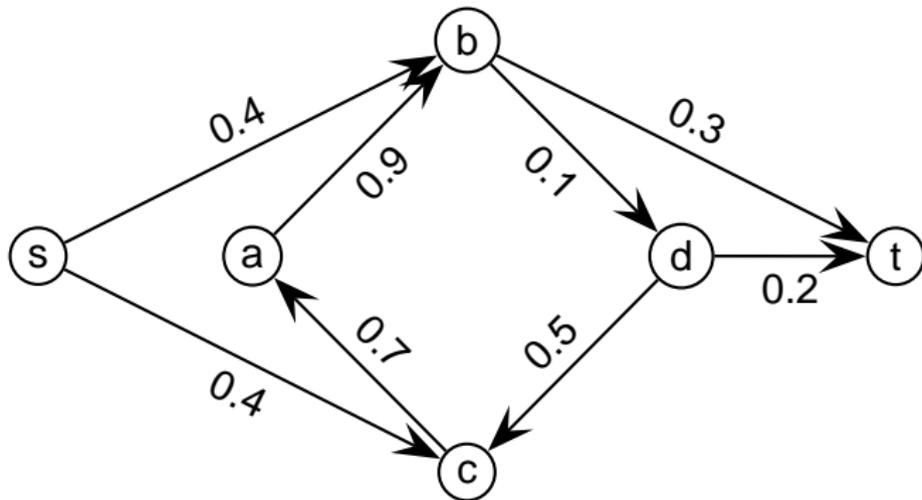
- $|\cdot|$ can be ∞ for infinite sets X or Y
- from now on only fuzzy relations on finite sets

Definition

Definition (Kawahara 2006)

An *s-t-network* N is a triple $N = (\alpha : X \leftrightarrow X, s, t)$, where α is a fuzzy endorelation on X with $\alpha \sqcap \alpha^\# = 0_{XX}$ and s (the *source*) and t (the *sink*) are two distinct elements of X .

Example s-t-network



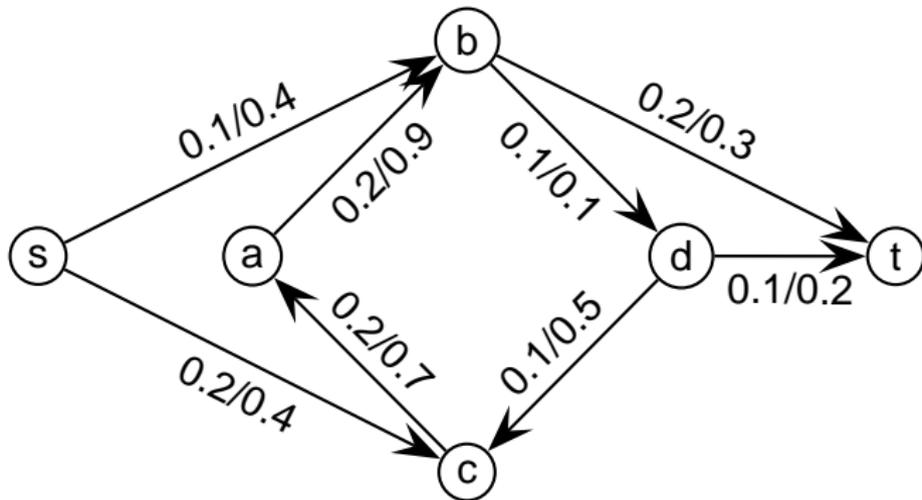
Definition

Definition

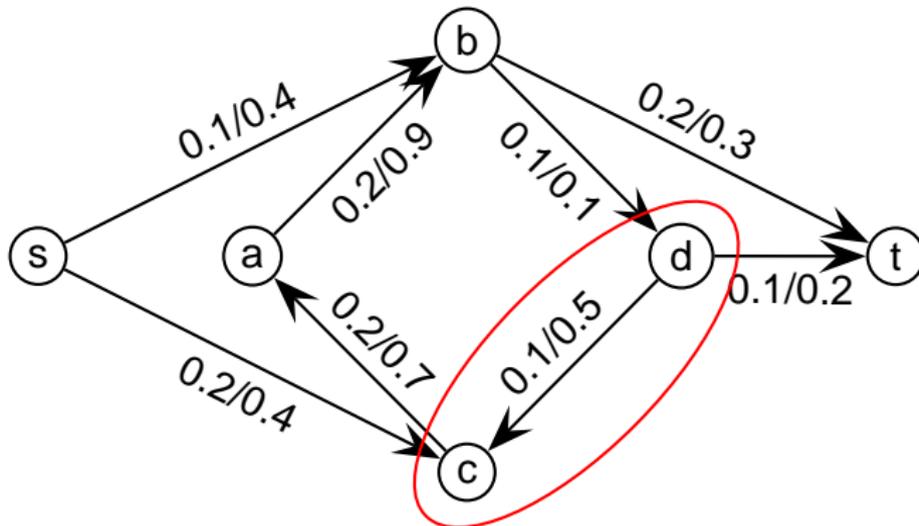
A *flow* φ on an s-t-network $N = (\alpha : X \leftrightarrow X, s, t)$ is a fuzzy endorelation on X with the properties

- $\varphi \sqsubseteq \alpha$ (capacity constraint)
- $|\tau\varphi| = |\varphi\tau|$ for all test relations τ on X with $\tau \sqcap (s \sqcup t) = 0_{XX}$ (flow conservation)

Example Flow



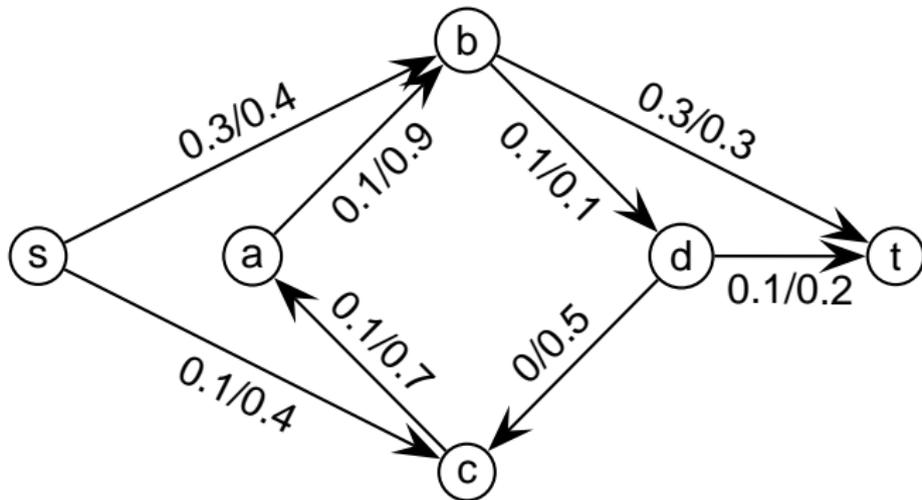
Example Flow



Value and Cut

- The *value* $val(\varphi)$ of a flow φ is defined as
$$val(\varphi) = |s\varphi| - |\varphi s|$$
- A flow is called *maximal*, if its value is maximal

Maximal Flow



Definition

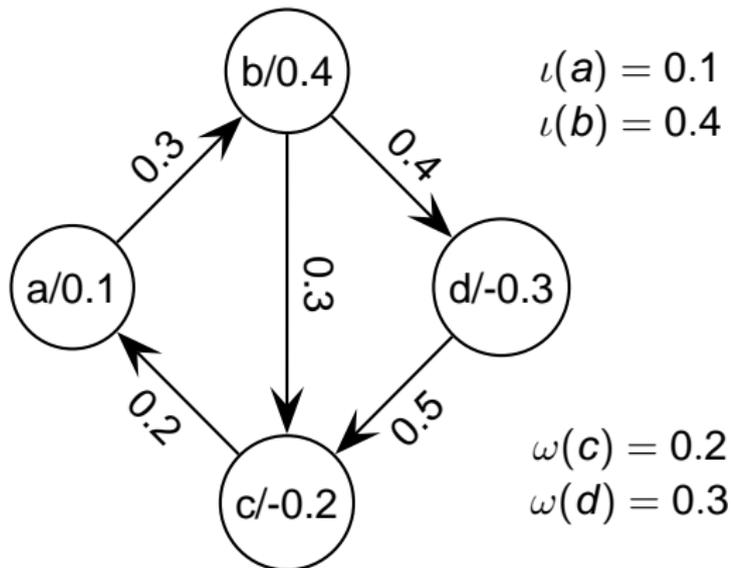
Definition

An *import network* I is a triple $I = (\alpha : X \leftrightarrow X, \iota, \omega)$ with the properties

- $\alpha \sqsubseteq 0.5 \cdot \nabla_{XX}$
- $\alpha \sqcap \alpha^\# = 0_{XX}$
- $\iota, \omega \sqsubseteq 0.5 \cdot id_X$
- $\iota \sqcap \omega = 0_{XX}$

α is called the *capacity function*, ι the *import* and ω the *outflow*.

Visualisation of Import Networks



Circulations on Import Networks

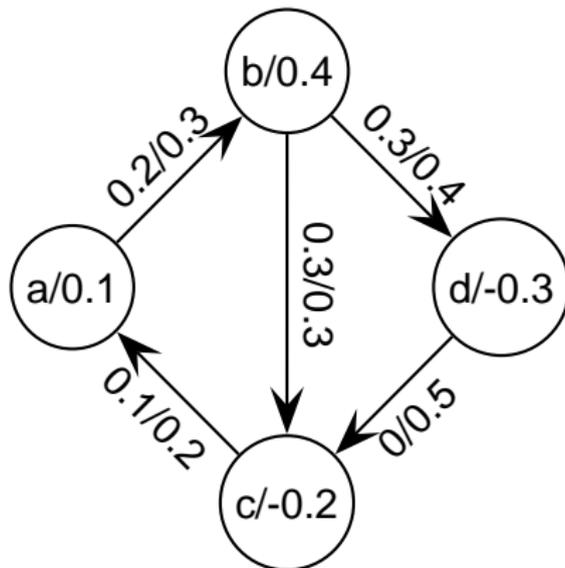
Definition

A *Circulation* φ on an import network $I = (\alpha : X \leftrightarrow X, \iota, \omega)$ is a fuzzy endorelation on X with the properties

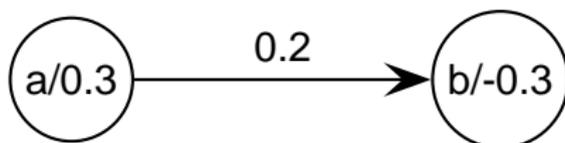
- $\varphi \sqsubseteq \alpha$ (capacity constraint)
- $|(\varphi \oplus \iota)\tau| = |\tau(\varphi \oplus \omega)|$ (flow conservation)

where $(\alpha \oplus \beta)(x, y) := \min\{1, \alpha(x, y) + \beta(x, y)\}$.

Example Circulation



Import Network without Circulation



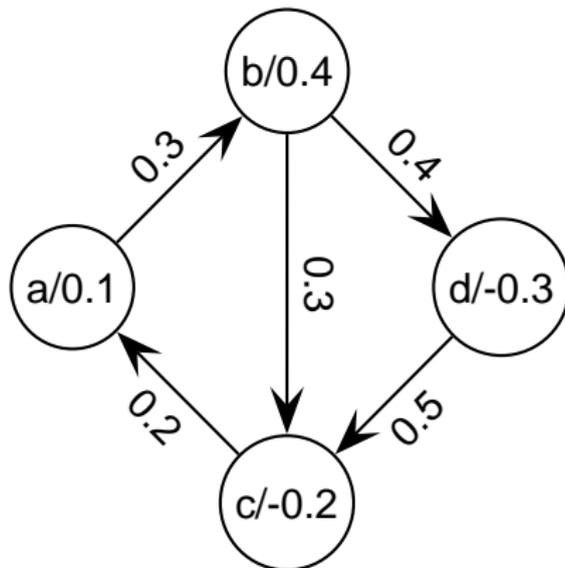
Existence and Determining of Circulations

- Existence of a circulation
- Determining a circulation

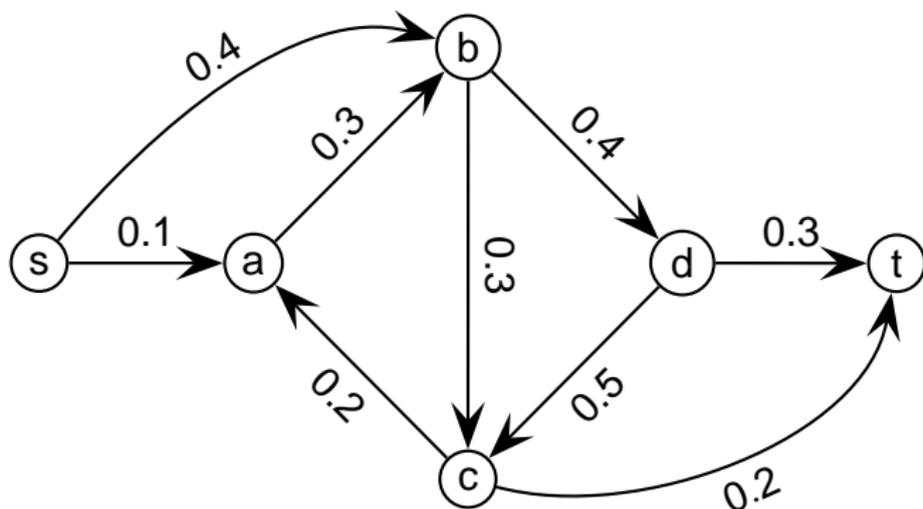
Existence and Determining of Circulations

- Existence of a circulation
- Determining a circulation
- Both via an s-t-network

Example Import Network



Associated s-t-Network

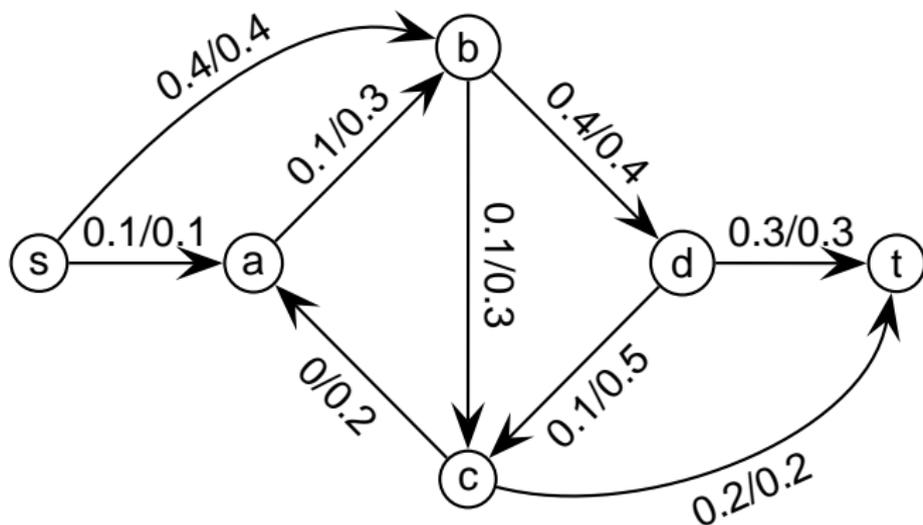


Existence Criterion for Circulations

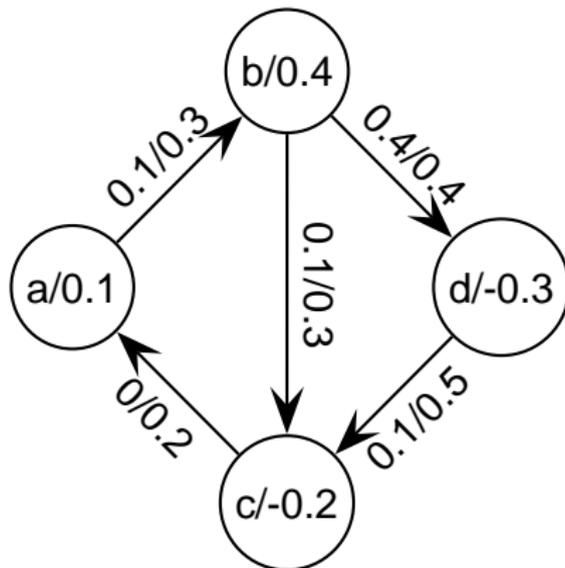
Theorem

There is a circulation on an import network iff the value of a maximal flow in the associated s-t-network equals both $|\iota|$ and $|\omega|$.

Maximal Flow in the Associated s-t-Network



Resulting Circulation



Summary

We saw

- an algebraic description of network problems
- application Kleene algebra in action
- first steps towards automated proving of theorems concerning networks

To Do's

- feeding automated provers with this stuff
- applying similar methods to related problems (partially done, MPC08)
- describing graph theoretical problems with these methods

Enjoy Your Meal!

Questions?
Criticisms?
Suggestions?